

## Minimization of FSA

Data Structures and Algorithms for Computational Linguistics III  
(ISCL-BA-07)

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## DFA minimization

- For any regular language, there is a unique *minimal* DFA
- By finding the minimal DFA, we can also prove equivalence (or not) of different FSA and the languages they recognize
- In general the idea is:
  - Throw away unreachable states (easy)
  - Merge equivalent states
- There are two well-known algorithms for minimization:
  - Hopcroft's algorithm: find and eliminate equivalent states by partitioning the set of states
  - Brzozowski's algorithm: "double reversal"

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## Finding equivalent states

Intuition



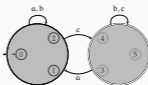
The edges leaving the group of nodes are identical.  
Their *right languages* are the same.

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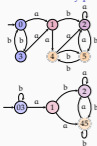


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## Minimization by partitioning



- Accepting & non-accepting states form a partition  
 $Q_1 = \{0, 1, 2, 3\}$ ,  $Q_2 = \{4, 5\}$
- If any two nodes go to different sets for any of the symbols split
- $Q_1 = \{0, 3\}$ ,  $Q_2 = \{1\}$ ,  $Q_3 = \{2, 4, 5\}$
- Stop when we cannot split any of the sets, merge the indistinguishable states

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## Minimization by partitioning

tabular version



- Create a state-by-state table, mark distinguishable pairs:  $(q_1, q_2)$  such that  $\Delta(q_1, x) \neq \Delta(q_2, x)$  is a distinguishable pair for any  $x \in \Sigma$



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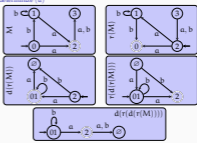
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- Merge indistinguishable states
- The algorithm can be improved by choosing which cell to visit carefully

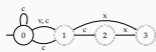
## Brzozowski's algorithm

double reverse (r), determinize (d)



## An exercise

find the minimum DFA for the automaton below



## Minimization algorithms

final remarks

- There are many versions of the 'partitioning' algorithm. General idea is to form equivalence classes based on right-language of each state.
- Partitioning algorithm has  $O(n \log n)$  complexity
- 'Double reversal' algorithm has exponential worst-time complexity
- Double reversal algorithm can also be used with NFAs (resulting in the minimal equivalent DFA - NFA minimization is intractable)
- In practice, there is no clear winner, different algorithms run faster on different input
- Reading suggestion: [hopcroft1979](#), [jurfafsky2009](#)

Next:

- PST
- PSA and regular languages

## Acknowledgments, credits, references

