

Bottom-up Chart Parsing: the CKY algorithm

Data Structures and Algorithms for Computational Linguistics III (ISCL-BA-07)

Çağrı Çöltekin
ccoltekin@ifa.uni-tuebingen.de

University of Tübingen
Seminar für Sprachwissenschaft

Winter Semester 2022/23

https://ifa.uni-tuebingen.de

Parsing so far

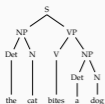
- Parsing is the task of automatic syntactic analysis
- For most practical purposes, context-free grammars are the most useful formalism for parsing
- We can formulate parsing as
 - Top-down: begin with the start symbol, try to *produce* the input string to be parsed
 - Bottom up: begin with the input, and try to *reduce* it to the start symbol
- Both strategies can be cast as search with backtracking
- Backtracking parsers are inefficient: they recompute sub-trees multiple times

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 1 / 13

Bottom-up parsing as search

Introduction CSF 1301



```
S → NP VP
NP → Det N
VP → V NP
VP → V
Det → a
Det → the
N → cat
N → dog
V → bites
N → bites
```

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 2 / 13

Dealing with ambiguity

Introduction CSF 1301

I saw her duck

```
S → NP VP
NP → NP N
NP → Prn
VP → V NP
VP → V
N → duck
V → duck
V → saw
Prn → I
Prn → she
Prn → her
```

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 3 / 13

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 4 / 13

Dealing with ambiguity

Introduction CSF 1301



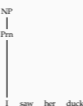
```
S → NP VP
NP → Prn N
NP → Prn
VP → V NP
VP → V
N → duck
V → duck
V → saw
Prn → I
Prn → she
Prn → her
```

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 5 / 13

Dealing with ambiguity

Introduction CSF 1301



```
S → NP VP
NP → Prn N
NP → Prn
VP → V NP
VP → V
N → duck
V → duck
V → saw
Prn → I
Prn → she
Prn → her
```

C. Çöltekin, MSR | University of Tübingen

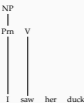
Winter Semester 2022/23 6 / 13

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 7 / 13

Dealing with ambiguity

Introduction CSF 1301



```
S → NP VP
NP → Prn N
NP → Prn
VP → V NP
VP → V
N → duck
V → duck
V → saw
Prn → I
Prn → she
Prn → her
```

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 8 / 13

Dealing with ambiguity

Introduction CSF 1301



```
S → NP VP
NP → Prn N
NP → Prn
VP → V NP
VP → V
N → duck
V → duck
V → saw
Prn → I
Prn → she
Prn → her
```

C. Çöltekin, MSR | University of Tübingen

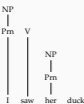
Winter Semester 2022/23 9 / 13

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 10 / 13

Dealing with ambiguity

Introduction CSF 1301



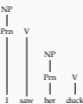
```
S → NP VP
NP → Prn N
NP → Prn
VP → V NP
VP → V
N → duck
V → duck
V → saw
Prn → I
Prn → she
Prn → her
```

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 11 / 13

Dealing with ambiguity

Introduction CSF 1301



```
S → NP VP
NP → Prn N
NP → Prn
VP → V NP
VP → V
N → duck
V → duck
V → saw
Prn → I
Prn → she
Prn → her
```

C. Çöltekin, MSR | University of Tübingen

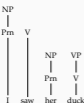
Winter Semester 2022/23 12 / 13

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 13 / 13

Dealing with ambiguity

Introduction CSF 1301



```
S → NP VP
NP → Prn N
NP → Prn
VP → V NP
VP → V
N → duck
V → duck
V → saw
Prn → I
Prn → she
Prn → her
```

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 14 / 13

Dealing with ambiguity

Introduction CSF 1301



```
S → NP VP
NP → Prn N
NP → Prn
VP → V NP
VP → V
N → duck
V → duck
V → saw
Prn → I
Prn → she
Prn → her
```

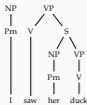
C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 15 / 13

C. Çöltekin, MSR | University of Tübingen

Winter Semester 2022/23 16 / 13

Dealing with ambiguity



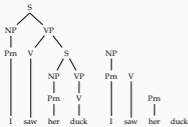
S → NP VP ←
 NP → Prn N
 NP → Prn
 VP → V NP
 VP → V
 VP → V S
 N → duck
 V → duck
 V → saw
 Prn → I
 Prn → she
 Prn → her

Dealing with ambiguity



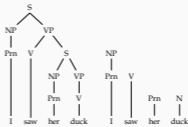
S → NP VP
 NP → Prn N
 NP → Prn
 VP → V NP
 VP → V
 VP → V S
 N → duck
 V → duck
 V → saw
 Prn → I
 Prn → she
 Prn → her

Dealing with ambiguity



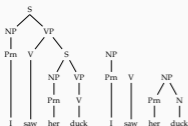
S → NP VP
 NP → Prn N
 NP → Prn
 VP → V NP
 VP → V
 VP → V S
 N → duck ←
 N → duck
 V → saw
 Prn → I
 Prn → she
 Prn → her

Dealing with ambiguity



S → NP VP
 NP → Prn N ←
 NP → Prn
 VP → V NP
 VP → V
 VP → V S
 N → duck
 V → duck
 V → saw
 Prn → I
 Prn → she
 Prn → her

Dealing with ambiguity



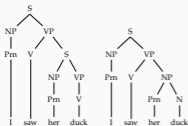
S → NP VP
 NP → Prn N
 NP → Prn
 VP → V NP ←
 VP → V
 VP → V S
 N → duck
 V → duck
 V → saw
 Prn → I
 Prn → she
 Prn → her

Dealing with ambiguity



S → NP VP ←
 NP → Prn N
 NP → Prn
 VP → V NP
 VP → V
 VP → V S
 N → duck
 V → duck
 V → saw
 Prn → I
 Prn → she
 Prn → her

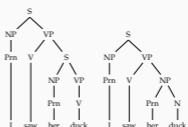
Dealing with ambiguity



S → NP VP
 NP → Prn N
 NP → Prn
 VP → V NP
 VP → V
 VP → V S
 N → duck
 V → duck
 V → saw
 Prn → I
 Prn → she
 Prn → her

How to represent multiple parses

parse forest grammar



$S_{2,4} \rightarrow NP_{0,1} VP_{1,4}$
 $NP_{2,3} \rightarrow Prn_{2,3}$
 $Prn_{0,1} \rightarrow I_{0,1}$
 $VP_{1,4} \rightarrow V_{1,2} S_{2,4}$
 $V_{1,2} \rightarrow saw_{1,2}$
 $S_{2,4} \rightarrow Prn_{2,3} V_{3,4}$
 $V_{3,4} \rightarrow duck_{3,4}$
 $VP_{1,4} \rightarrow V_{1,2} NP_{2,4}$
 $NP_{2,4} \rightarrow Prn_{2,3} N_{3,4}$

CKY algorithm

- The CKY (Cocke-Kasami-Younger) parsing algorithm is a dynamic programming algorithm
- It processes the input *bottom up*, and saves the intermediate results on a chart
- Time complexity for recognition is $O(n^3)$
- Space complexity is $O(n^2)$
- It requires the CFG to be in Chomsky normal form (CNF) (can somewhat be relaxed, but not common)

Chomsky normal form (CNF)

- A CFG is in CNF, if the rewrite rules are in one of the following forms
 - $A \rightarrow BC$
 - $A \rightarrow a$
 where A, B, C are non-terminals and a is a terminal
- Any CFG can be converted to CNF
- Resulting grammar is *not* equivalent to the original grammar:
 - it generates/accepts the same language
 - but the derivations are different

Converting to CNF: example

S → NP VP
 S → Aux NP VP
 NP → the N
 VP → V NP
 VP → V
 N → cat
 N → dog
 V → bites
 N → bites

• S → Aux NP VP
 S → Aux NP VP → S → Aux X
 X → NP VP

• NP → the N
 NP → the N → NP → X N
 X → the

• VP → V
 VP → V → VP → bites

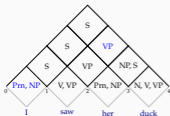
Converting to CNF

- Eliminate the ϵ rules: if $A \rightarrow \epsilon$ is in the grammar
 - replace any rule $B \rightarrow \alpha A \beta$ with two rules
 - $B \rightarrow \alpha \beta$
 - $B \rightarrow \alpha A' \beta$
 - add $A' \rightarrow \alpha$ for all α (except ϵ) whose LHS is A
 - repeat the process for newly created ϵ rules
 - remove the rules with ϵ on the RHS (except $S \rightarrow \epsilon$)
- Eliminate unit rules: for a rule $A \rightarrow B$
 - Replace the rule with $A \rightarrow a_1 | \dots | a_n$, where a_1, \dots, a_n are all RHS of rule B
 - Remove the rule $A \rightarrow B$
 - Repeat the process until no unit rules remain
- Binarize all the non-binary rules with non-terminal on the RHS: for a rule $A \rightarrow X_1 X_2 \dots X_n$:
 - Replace the rule with $A \rightarrow A_1 X_1 \dots X_n$, and add $A_1 \rightarrow X_1 X_2$
 - Repeat the process until all new rules are binary

CKY demonstration

an ambiguous example

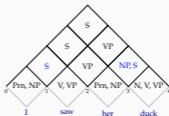
S → NP VP



S → NP VP
 NP → Prn N
 VP → V NP
 VP → V S
 N → duck
 VP → duck | saw
 V → duck | saw
 Prn → I | she | her
 NP → I | she | her

CKY demonstration

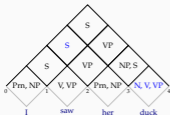
an ambiguous example



S → NP VP
 NP → Prn N
 VP → V NP
 VP → V S
 N → duck
 VP → duck | saw
 V → duck | saw
 Prn → I | she | her
 NP → I | she | her

CKY demonstration

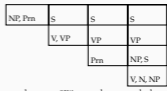
an ambiguous example



S → NP VP
 NP → Prn N
 VP → V NP
 VP → V S
 N → duck
 VP → duck | saw
 V → duck | saw
 Prn → I | she | her
 NP → I | she | her

CKY demonstration: the chart

our chart is a 2D array

Space complexity is $O(n^2)$.

CKY demonstration: the chart

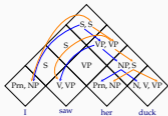
our chart is a 2D array – this is more convenient for programming

Space complexity is $O(n^2)$.

Parsing vs. recognition

- We went through a recognition example
- Note that the algorithm is not directional: it takes the complete input
- Recognition accepts or rejects a sentence based on a grammar
- For parsing, we want to know the derivations that yielded a correct parse
- To recover parse trees, we
 - follow the same procedure as recognition
 - add back links to keep track of the derivations

Chart parsing example (CKY parsing)



The chart stores a parse forest efficiently.

Summary

- + CKY avoids re-computing the analyses by storing the earlier analyses (of sub-spans) in a table
- It still computes lower level constituents that are not allowed by the grammar
- CKY requires the grammar to be in CNF
- CKY has $O(n^3)$ recognition complexity
- For parsing, we need to keep track of backlinks
- CKY can efficiently store all possible parses in a chart
- Enumerating all possible parses have exponential complexity (worst case)
- Suggested reading: **Jurafsky 2009**

Next:

- Top-down chart parsing: Earley algorithm
- Suggested reading:
 - **Jurafsky 2009**
 - **grune 2008**

Acknowledgments, references, additional reading material